

## CS122 Algorithms and Data Structures

MW 11:00 am - 12:15 pm, MSEC 101

Instructor: Xiao Qin

Lecture 9: Binary Trees

## Why Trees?

### n Limitations of

- Arrays
- Linked lists
- Stacks
- Queues

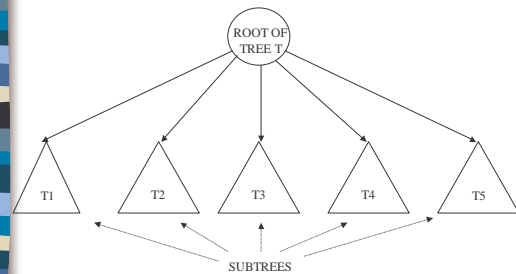


## Trees: Recursive Definition

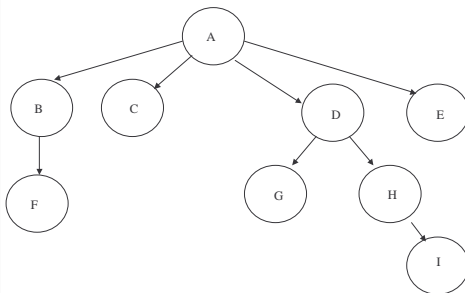


- n A tree is a collection of nodes.
- n The collection can be empty, or consist of a “root” node  $R$ .
- n There is a “directed edge” from  $R$  to the root of each subtree. The root of each subtree is a “child” of  $R$ .  $R$  is the “parent” of each subtree root.

## Trees: Recursive Definition (cont.)



## Trees: An Example



## Trees: More Definitions

- n Nodes with no children are **leaves**: (C,E,F,H,I).
- n Nodes with the same parents are **siblings**: (B,C,D,E) and (G,H).
- n A **path** from node  $n$  to node  $m$  is the sequence of directed edges from  $n$  to  $m$ .
- n A **length of a path** is the number of edges in the path

## Trees: More Definitions (cont.)

- n The **level/depth of node  $n$**  is the length of the path from the root to  $n$ . The level of the root is 0.
- n The **height/depth of a tree** is equal to the maximum level of a node in the tree.
- n The **height of a node  $n$**  is the length of the longest path from  $n$  to a leaf. The height of a leaf node is 0.
- n The height of a tree is equal to the height of the root node.

## Binary Trees – A Informal Definition

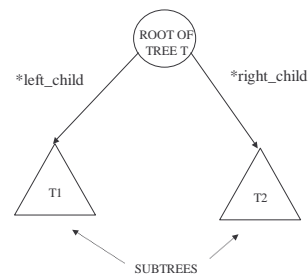
- n A binary tree is a tree in which no node can have more than two children.
- n Each node has 0, 1, or 2 children
  - In this case we can keep direct links to the children:

```
struct TreeNode
{
    Object element;
    TreeNode *left_child;
    TreeNode *right_child;
};
```

## Binary Trees – Formal Definition

- n A binary tree is a structure that
  - contains no nodes, or
  - is comprised of three disjoint sets of nodes:
    - a **root**
    - a binary tree called its **left subtree**
    - a binary tree called its **right subtree**
- n A binary tree that contains no nodes is called **empty**

## Binary Trees: Recursive Definition



## Differences Between A Tree & A Binary Tree

- n No node in a binary tree may have more than 2 children, whereas there is no limit on the number of children of a node in a tree.
- n The subtrees of a binary tree are ordered; those of a tree are not ordered.

## Differences Between A Tree & A Binary Tree (cont.)

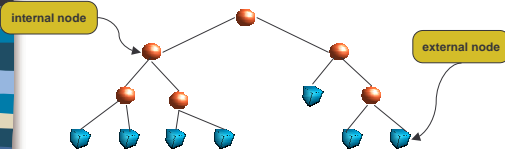
- n The subtrees of a binary tree are ordered; those of a tree are not ordered



- Are different when viewed as binary trees
- Are the same when viewed as trees

## Internal and External Nodes

- n Because in a binary tree all the nodes must have the same number of children we are forced to change the concepts slightly
  - We say that all **internal nodes** have two children
  - **External nodes** have no children

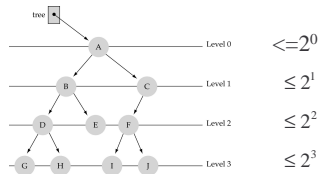


## Recursive definition of a Binary Tree

- n Most of concepts related to binary trees can be explained recursive
- n For instance, A binary tree is:
  - An **external node**, or
  - An **internal node** connected to a left binary tree and a right binary tree (called left and right subtrees)
- n In programming terms we can see that our definition for a linked list (singly) can be modified to have two links from each node instead of one.

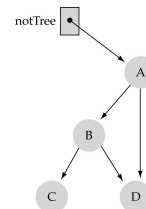
## What is a binary tree?

- 1 **Property1:** each node can have up to two successor nodes (**children**)
  - 1 The predecessor node of a node is called its **parent**
  - 1 The "beginning" node is called the **root** (no parent)
  - 1 A node without **children** is called a **leaf**



## What is a binary tree? (cont.)

- n **Property2:** a unique path exists from the root to every other node

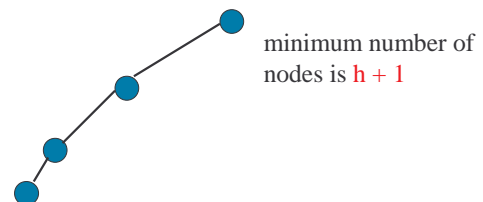


## Mathematical Properties of Binary Trees

- n Let's us look at some important mathematical properties of binary trees
- n A good understanding of these properties will help the understanding of the performance of algorithms that process trees
- n Some of the properties we'll describe relate also the structural properties of these trees. This is the case because performance characteristics of many algorithms depend on these structural properties and not only the number of nodes.

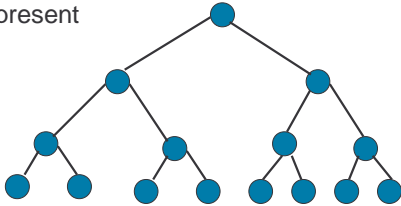
## Minimum Number Of Nodes

- n Minimum number of nodes in a binary tree whose height is **h**.
- n At least one node at each level.



## Maximum Number Of Nodes

- n All possible nodes at first  $h$  levels are present



Maximum number of nodes

$$= 1 + 2 + 4 + 8 + \dots + 2^h = 2^{h+1} - 1$$

## Number Of Nodes & Height

- n Let  $n$  be the number of nodes in a binary tree whose height is  $h$ .
- n  $h + 1 \leq n \leq 2^{h+1} - 1$
- n  $\log_2(n+1) - 1 \leq h \leq n - 1$
- n The max height of a tree with  $N$  nodes is  $N - 1$  (same as a linked list)
- n The min height of a tree with  $N$  nodes is  $\log(N+1) - 1$

## Relationship Between Number of Nodes (Internal - External)

- n A binary tree with  $N$  internal nodes has  $N+1$  external nodes

*Let's try to prove this using induction...*

## Number of edges

- n A binary tree with  $N$  internal nodes has  $2N$  edges

*Let's try to prove this using induction...*

## Number of edges

- n A binary tree with  $N$  nodes (internal and external) has  $N-1$  edges

*Let's try to prove this using induction...*

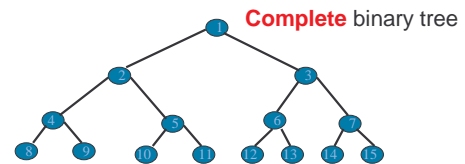
## Binary Tree Representation

- n Array representation
- n Linked representation

## Binary Trees

- n **Full** binary tree :
  - All **internal** nodes have two children.
- n **Complete** binary tree :
  - All leaves have the same level
  - All **internal nodes** have two children

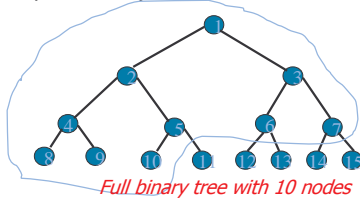
## Node Number Properties



- n Parent of node  $i$  is node  $i/2$ 
  - But node 1 is the root and has no parent
- n Left child of node  $i$  is node  $2i$ 
  - But if  $2i > n$ , node  $i$  has no left child
- n Right child of node  $i$  is node  $2i+1$ 
  - But if  $2i+1 > n$ , node  $i$  has no right child

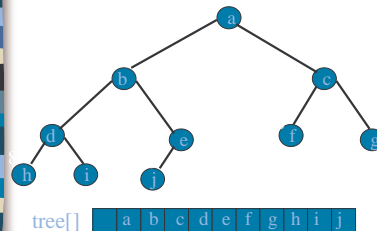
## Full Binary Tree With $n$ Nodes

- n Start with a full binary tree that has at least  $n$  nodes.
- n Number the nodes as described earlier.
- n The binary tree defined by the nodes numbered 1 through  $n$  is the unique  $n$  node complete binary tree.



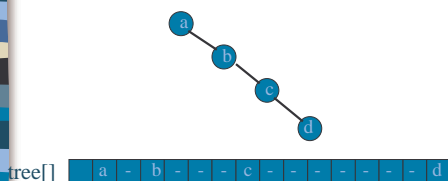
## Array Representation

- n Number the nodes using the numbering scheme for a full binary tree
- n Store the node numbered  $i$  in `tree[i]`



## Right-Skewed Binary Tree

- n An  $n$  node binary tree needs an array whose length is between  $n+1$  and  $2n$



## Linked Representation

- n Each tree node is represented as an object whose data type is `TreeNode`
- n The space required by an  $n$  node binary tree is  $n * (\text{space required by one node})$

## Trees: Linked representation Implementation 1

```
struct TreeNode
{
    Object element;
    TreeNode *child1;
    TreeNode *child2;
    .
    .
    .
    TreeNode *childn;
};
```

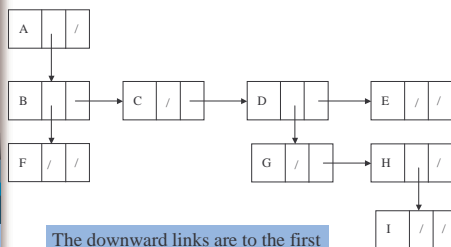
§ Each node contains a link to all of its children.  
§ This isn't a good idea, because a node can have an arbitrary number of children!

## Trees: Linked representation Implementation 2

```
struct TreeNode
{
    Object element;
    TreeNode *child1;
    TreeNode *sibling;
};
```

Each node contains a link to its first child and a link to its next sibling. This is a better idea.

## Implementation 2: Example



The downward links are to the first child; the horizontal links are to the next sibling.

## Binary Trees

n A **binary tree** is a tree whose nodes have at most two offspring

n **Example**

```
struct nodeType {
    object element;
    struct nodeType *left, *right;
};
struct nodeType *tree;
```

## Linked Representation Example

