

CS122 Algorithms and Data Structures

MW 11:00 am - 12:15 pm, MSEC 101

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Lecture 5: Doubly Linked Lists and
Linked Stacks

Lists

- A list is a varying-length, **linear collection** of homogeneous elements.
- Linear means each list element (except the first) has a **unique predecessor**, and each element (except the last) has a **unique successor**.
- In a linear structure, components can only be **accessed sequentially** one after the other.

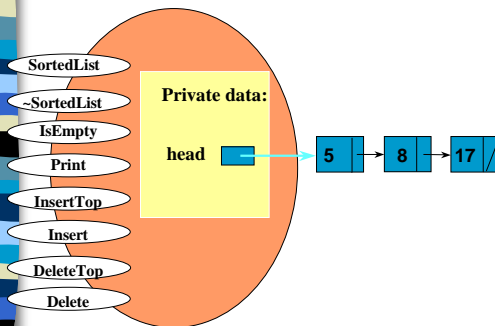
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An Insert Algorithm for Sorted Lists

- **Sorted list:** list elements sorted in ascending order
- How would the algorithm to insert an item into a sorted linked list?

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class SortedList



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ADT SortedList Operations

Transformers
– InsertTop
– Insert
– DeleteTop
– Delete



Observers
– Print
– IsEmpty



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struct NodeType

// SPECIFICATION FILE DYNAMIC-LINKED SORTED LIST
(slist.h)

typedef int ItemType; // Type of each component
// is simple type or string type

```
struct NodeType
{
    ItemType item; // Pointer to person's name
    NodeType* next; // link to next node in list
};
```

typedef NodeType* NodePtr;

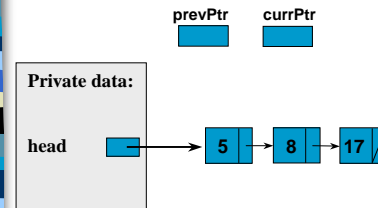
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Insert algorithm for SortedList

- find proper position for the new element in the sorted list using **two pointers prevPtr and currPtr**, where prevPtr trails behind currPtr
- obtain a node for insertion and place item in it
- **insert the node by adjusting pointers**

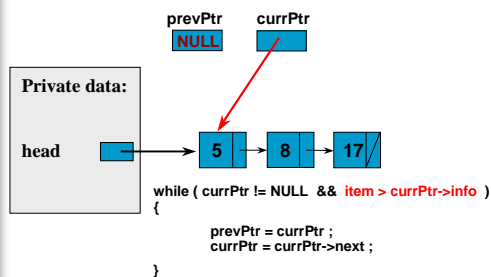
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Inserting 12 into a Sorted List



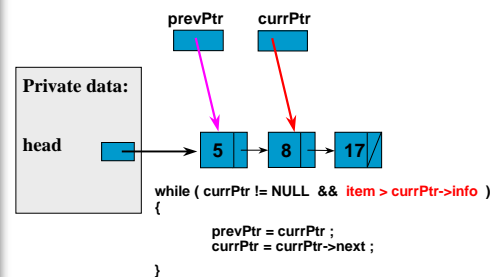
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Finding Proper Position for 12



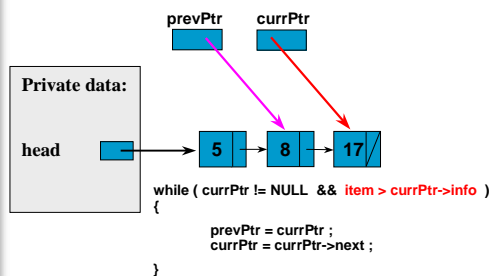
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Finding Proper Position for 12



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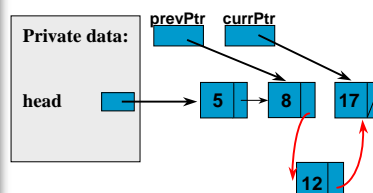
Finding Proper Position for 12



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Inserting '12' into Proper Position

```
location->next = currPtr ;
if ( prevPtr == NULL )
    head = location ;
else
    prevPtr->next = location ;
```



12

// IMPLEMENTATION DYNAMIC-LINKED SORTED LIST // (slist.cpp)

```
SortedList::SortedList () // Constructor
// Post: head == NULL
{
    head = NULL ;
}

SortedList::~~SortedList2 () // Destructor
// Post: All linked nodes deallocated
{
    ItemType temp ;
    // keep deleting top node
    while ( !IsEmpty )
        DeleteTop ( temp );
}
```

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```
void SortedList::Insert( ItemType item )
// Pre: item is assigned && list components in ascending order
// Post: new node containing item is in its proper place
// && list components in ascending order
{
    NodePtr currPtr, prevPtr, location ;
    location = new NodeType ;
    newNodePtr->info = item ;
    prevPtr = NULL ;
    currPtr = head ;
    while ( currPtr != NULL && item > currPtr->info )
    {
        prevPtr = currPtr ; // advance both pointers
        currPtr = currPtr->next ;
    }
    location->next = currPtr ; // insert new node here
    if ( prevPtr == NULL )
        head = location ;
    else
        prevPtr->next = location ;
}
```

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```
void SortedList::DeleteTop ( ItemType& item )
// Pre: list is not empty && list elements in ascending order
// Post: item == element of first list node @ entry
// && node containing item is no longer in linked list
// && list elements in ascending order
{
    NodePtr tempPtr = head ;

    item = head->info ; // obtain item
    head = head->next ; // advance head
    delete tempPtr ;
}
```

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```
void SortedList::Delete ( /* in */ ItemType item )
// Pre: list is not empty && list elements in ascending order
// && item == component member of some list node
// Post: item == element of first list node @ entry
// && node containing first occurrence of item is no longer
// in linked list && list elements in ascending order
{
    NodePtr delPtr ;
    NodePtr currPtr ; // Is item in first node?
    if ( item == head->info )
    {
        delPtr = head ; // If so, delete first node
        head = head->next ;
    }
    else { // search for item in rest of list
        currPtr = head ;
        while ( currPtr->next->info != item )
            currPtr = currPtr->next ;
        delPtr = currPtr->next ;
        currPtr->next = currPtr->next->next ;
    }
    delete delPtr ;
}
```

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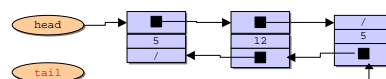
Doubly Linked Lists

- An extension of a Singly Linked List
- Each node has two pointer
 - One pointing to the successor
 - One pointing to the predecessor
- They are used because they ease certain operations like the deleteElement
- They are interesting for traversal as you can move in either directions

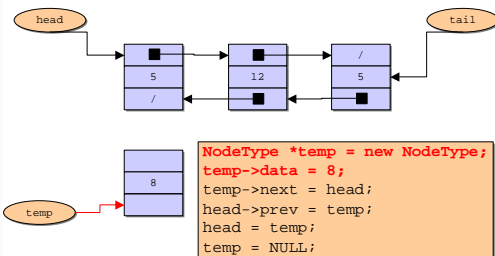
```
struct NodeType {
    <someType> data;
    NodeType* prev;
    NodeType* next;
}
```

Doubly Linked Lists (Cont.)

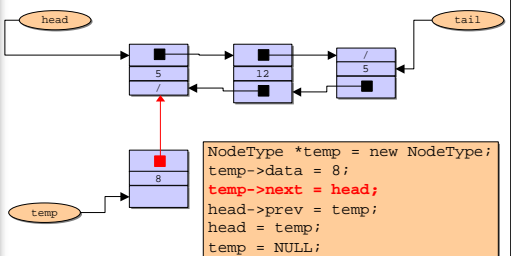
- Most of our structures and algorithms will be implemented with the help of Doubly Linked Lists
- Although some operations are made easier to understand they also become a bit slower due to the overhead of extra pointers
- In addition to the head a pointer called **tail** is also maintained



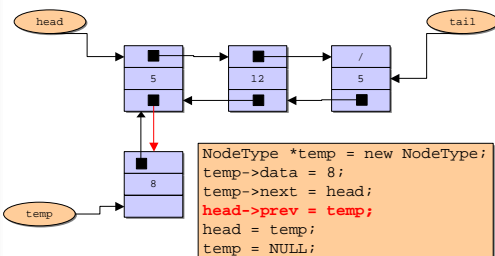
Inserting in the Front



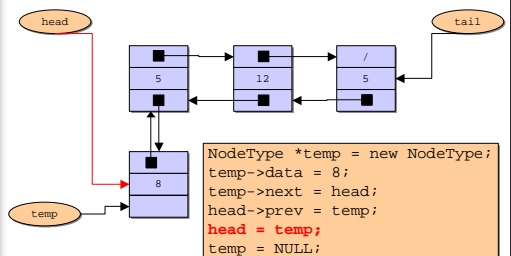
Inserting in the Front (cont.)



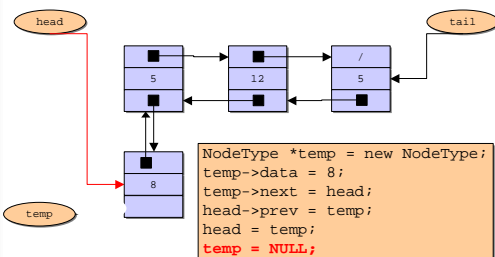
Inserting in the Front (cont.)



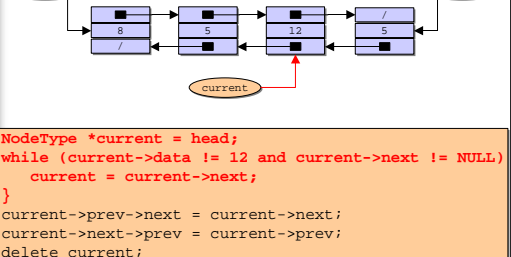
Inserting in the Front (cont.)



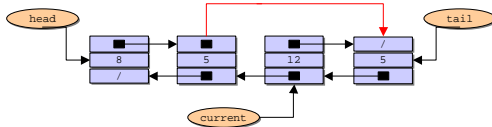
Inserting in the Front (cont.)



Deleting element '12'



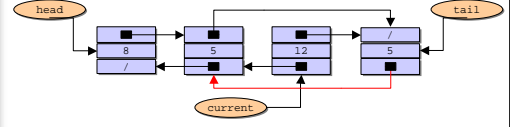
Deleting element '12' (cont.)



```

NodeType *current = head;
while (current->data != 12 and current->next != NULL)
    current = current->next;
}
current->prev->next = current->next;
current->next->prev = current->prev;
delete current;
    
```

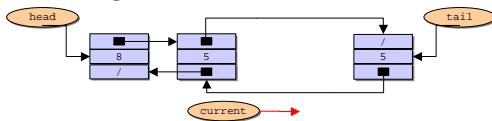
Deleting element '12' (cont.)



```

NodeType *current = head;
while (current->data != 12 and current->next != NULL)
    current = current->next;
}
current->prev->next = current->next;
current->next->prev = current->prev;
delete current;
    
```

Deleting element '12' (cont.)

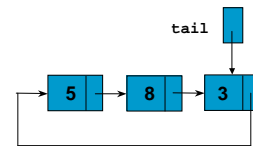


```

NodeType *current = head;
while (current->data != 12 and current->next != NULL)
    current = current->next;
}
current->prev->next = current->next;
current->next->prev = current->prev;
delete current;
    
```

Circular Lists

- Nodes form a ring
- Each node has a successor

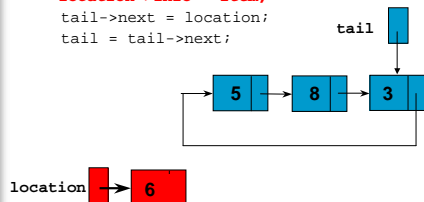


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Inserting a Node at the Tail of a List Circular Lists

```

int    item = 6;
Node  *location;
location = new NodeType;
location->info = item;
tail->next = location;
tail = tail->next;
    
```

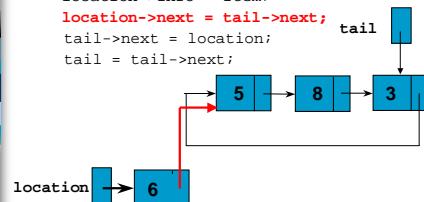


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Inserting a Node at the Tail of a List Circular Lists (cont.)

```

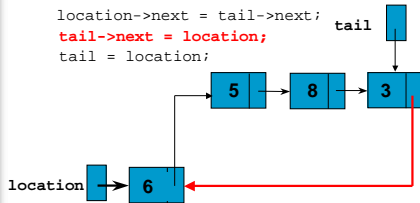
int    item = 6;
Node  *location;
location = new NodeType;
location->info = item;
location->next = tail->next;
tail->next = location;
tail = tail->next;
    
```



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Inserting a Node at the Tail of a List Circular Lists (cont.)

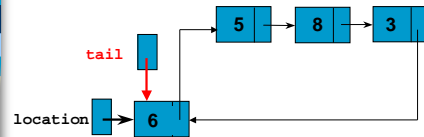
```
int    item = 6;
Node  *location;
location = new NodeType;
location->info = item;
location->next = tail->next;
tail->next = location;
tail = location;
```



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Inserting a Node at the Tail of a List Circular Lists (cont.)

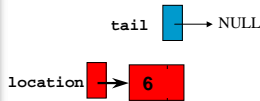
```
int    item = 6;
Node  *location;
location = new NodeType;
location->info = item;
location->next = tail->next;
tail->next = location;
tail = location;
```



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Inserting a Node at the Tail of a List Circular Lists: A Special Case

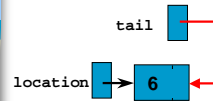
```
int    item = 6;
Node  *location;
location = new NodeType;
location->info = item;
if (isEmpty()) {
    tail = location;
    tail->next = tail;
}
```



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Inserting a Node at the Tail of a List Circular Lists: A Special Case (cont.)

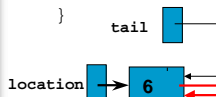
```
int    item = 6;
Node  *location;
location = new NodeType;
location->info = item;
if (isEmpty()) {
    tail = location;
    tail->next = tail;
}
```



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Inserting a Node at the Tail of a List Circular Lists: A Special Case (cont.)

```
int    item = 6;
Node  *location;
location = new NodeType;
location->info = item;
if (isEmpty()) {
    tail = location;
    tail->next = tail;
} else {
    .....
```



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Debugging Linked Lists

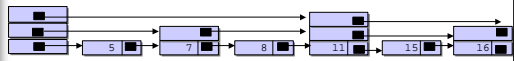
Two Suggestions Only

- Draw your linked list.
- Consider special cases.

Skip Lists

- Linked lists require sequential scanning for a search operation
- Skip lists allow for skipping certain nodes
- A skip list is an ordered linked list

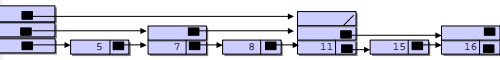
A Skip List with Evenly spaced Nodes of different levels



A Find Algorithm for Skip Lists

```

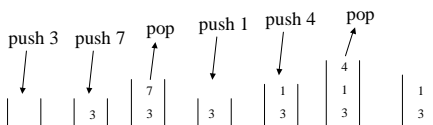
Find(element el)
p = the non-null list on the highest level i;
while el not found and level i >= 0
    if p->key > el //eg. el = 8
        p = a sublist that begins in the predecessor of p on level i--;
    else if p->key < el
        if p is the last node on level i; //eg. el = 16
            p = a nonnull sublist begins in p on the highest level < i;
            i = the number of the new level; //i = i-1
        else p = p->next;
    else return(TRUE);
return(FALSE);
    
```



Stacks using Linked Lists

- Implementation of stacks using linked lists are very simple
- The difference between a normal linked list and a stack using a linked list is that some of the linked list operations are not available for stacks
- Being a stack we have only one insert operation called `push()`.
 - In many ways push is the same as insert in the front
- We have also one delete operation called `pop()`
 - This operation is the same as the operation delete from the front
- The other important operations in a stack, called `top()` and `isEmpty()`, don't modify the structure
- To implement stacks we will use a singly linked list

A series of operations executed on a stack



Pictorial view of a stack

