

## CS122 Algorithms and Data Structures

MW 11:00 am - 12:15 pm, MSEC 101

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Lecture 19: Sorting (1)

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## Introduction

- n Common problem: sort a list of values, starting from lowest to highest.
  - Telephone directory
  - Words of dictionary in alphabetical order
  - Students names listed alphabetically
- n Choose a **criteria** which is used to order data
- n Given a list of records that have **keys**, we use these keys to define an ordering of the items in the list.

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## Elementary Sorting Algorithms

- n We are given  $n$  records to sort.
- n There are a number of simple sorting algorithms whose worst and average case performance is quadratic  $O(n^2)$ :
  - Insertion sort
  - Selection sort
  - Bubble sort

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## The Insertion Sort Algorithm

- n Given an array of integers
- n The Insertion Sort algorithm views the array as having a **sorted side** and an **unsorted side**.
- n The sorted side starts with just the first element, which is not necessarily the smallest element.
- n The sorted side grows by taking the front element from the unsorted side.

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## The Insertion Sort Algorithm (cont.)

- n The sorted side grows by **taking the front element** from the unsorted side.
- n **Inserting it** in the place that keeps the sorted side arranged from small to large.
- n In some cases there is **no need to move** the new inserted item.

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## How to Insert an Element

- n **Copy** the new element to a separate location.
- n **Shift** elements in the sorted side, creating an open space for the new element.
- n **Continue shifting** elements until you reach the location for the new element.
- n **Copy** the new element **back** into the array, at the correct location.

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## The Insertion Sort Algorithm

```
InsertionSort(data[], n)
  for i = 1 to n-1
    take next key from unsorted part of array
    insert in appropriate location in sorted part of array:
      for j = i down to 0,
        shift sorted elements to the right if data > data[j]
    increase size of sorted array by 1
```

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## Insertion Sort Time Analysis

- n In O-notation, what is:
  - Worst case running time for  $n$  items?
  - Average case running time for  $n$  items?

- n Steps of algorithm:

```
for i = 1 to n-1  Outer loop:  $O(n)$ 
  take next key from unsorted part of array
  insert in appropriate location in sorted part of array:
    for j = i down to 0,  Inner loop:  $O(n)$ 
      shift sorted elements to the right if key > key[j]
  increase size of sorted array by 1
```

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## The Selection Sort Algorithm

- n Start by finding the **smallest** element.
- n Swap the smallest entry with the **first element**.
- n Part of the array is sorted.
- n Find the smallest element in the unsorted side.
- n Swap with the front of the unsorted side.
- n The size of the sorted side is increased by one element.
- n **Continue** until the unsorted side has just one number. Why?

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## The Selection Sort Algorithm (cont.)

Basic Idea: Repeatedly **select** the smallest element, and **move** this element to the front of the unsorted side.

```
Selectsort(data[], n)
  for i = 1 to n-1
    find smallest key in unsorted part of array;
    swap smallest item to front of unsorted array;
    decrease size of unsorted array by 1;
```

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## Selection Time Sort Analysis

- n In O-notation, what is:
  - Worst case running time for  $n$  items?
  - Average case running time for  $n$  items?
- n Steps of algorithm:
  - for i = 1 to n-1  $O(n)$ 
    - find smallest key in unsorted part of array  $O(n)$
    - swap smallest item to front of unsorted array
    - decrease size of unsorted array by 1
- n Selection sort analysis:  $O(n^2)$

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## The Bubble Sort Algorithm

- n **Scan** the array from right to left.
- n Look at pairs of elements (**adjacent elements**) in the array, and swaps their order if needed.
- n **Repeatedly scan** the array from right to left elements until you reach the location for the new element.
- n Continue scanning until done

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## The Bubble Sort Algorithm (cont.)

```
BubbleSort(data[ ], n)
  for i = 0 to n-1
    for j = 0 to n-1
      if data[j] > data[j+1] then swap
    if no elements swapped in this pass through array, done.
  otherwise, continue
```

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## Bubble Sort Time Analysis

Steps of algorithm:

```
for i = 0 to n-1          O(n)
  for j = 0 to n-1        O(n)
    if data[j] > data[j+1] then swap
  if no elements swapped in this pass through array,
  done.
  otherwise, continue
```

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## Conclusions

- n Selection Sort, Insertion Sort, and Bubble Sort all have a worst-case time of  $O(n^2)$ , making them impractical for large arrays.
- n But they are easy to program, easy to debug.
- n Insertion Sort also has good performance when the array is nearly sorted to begin with.
- n But more sophisticated sorting algorithms are needed when good performance is needed in all cases for large arrays.
- n Next time: Quick Sort, Merge Sort, and Radix Sort.

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