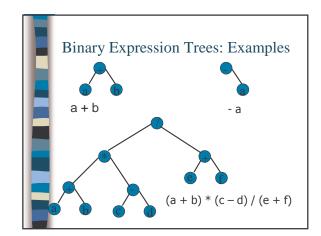
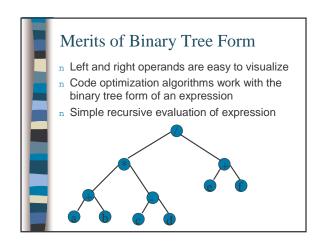
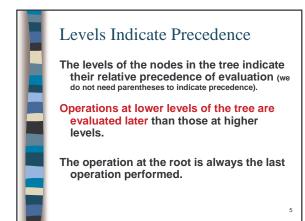
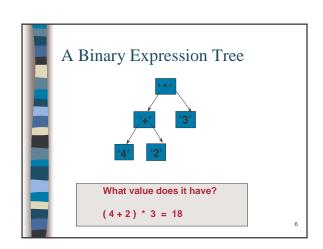
CS122 Algorithms and Data Structures MW 11:00 am - 12:15 pm, MSEC 101 Instructor: Xiao Qin Lecture 11: Binary Expression Trees

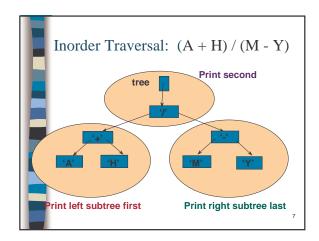


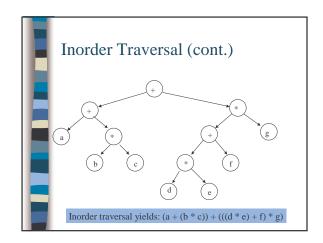


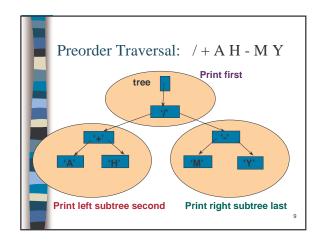


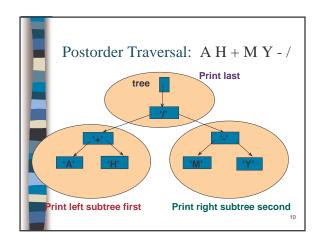




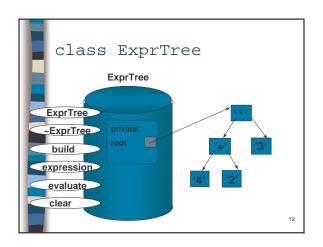






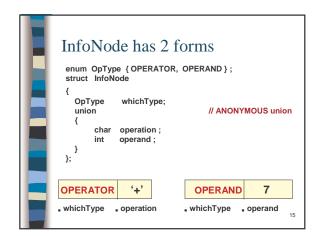


Traversals and Expressions In Inorder traversal produces the infix representation of the expression. In postorder traversal produces the postfix representation of the expression. In Preorder traversal produces a representation that is the same as the way that the programming language Lisp processes arithmetic expressions!



```
class ExprTree {
public:
  ExprTree ();
                     // Constructor
  ~ExprTree ();
                     // Destructor
  void build ();
                     // build tree from prefix expression
  void expression () const;
  // output expression in fully parenthesized infix form
  float evaluate () const;
                                  // evaluate expression
                                  // clear tree
  void clear ():
  void showStructure () const; // display tree
private:
  void showSub();
                          // recursive partners
  struct TreeNode *root;
};
```

```
Each node contains two pointers
struct TreeNode
  InfoNode info;
                     // Data member
  TreeNode* left;
                     // Pointer to left child
                     // Pointer to right child
  TreeNode* right;
  NULL
                               6000
                      7
            OPERAND
           whichType
                     . operand
  left
               info
                              right
```



```
int Eval(TreeNode* ptr)
{
    switch (ptr->info.whichType)
    {
        case OPERAND: return ptr->info.operand;
        case OPERATOR:
        switch (tree->info.operation)
        {
            case '+' : return (Eval(ptr->left) + Eval(ptr->right));
            case '-' : return (Eval(ptr->left) - Eval(ptr->right));
            case '*' : return (Eval(ptr->left) * Eval(ptr->right));
            case 'f' : return (Eval(ptr->left) / Eval(ptr->right));
        }
    }
}
```

Constructing an Expression Tree n There is a simple O(N) stack-based algorithm to convert a postfix expression into an expression tree. n Recall we also have an algorithm to convert an infix expression into postfix, so we can also convert an infix expression into an expression tree without difficulty (in O(N) time).

